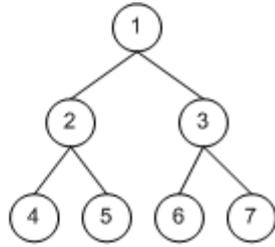


(a)



(b)

# CSSE 230 Day 10

## Size vs height in a Binary Tree

After today, you should be able to...

... use the relationship between the size and height of a tree to find the maximum and minimum number of nodes a binary tree can have

... understand the idea of mathematical induction as a proof technique

# Today

- ▶ New material:
  - Size vs height of trees: patterns and proofs
- ▶ Review for test next class
  - Written (50–70%):
    - big  $O/\theta/\Omega$ : true/false, using definitions, code analysis
    - Choosing an ADT to solve a given problem
    - Maybe a bit with binary trees
  - Programming (30–50%):
    - Implementing one ADT using another ADT
- ▶ Due after that:
  - Hardy's Taxi, part two: efficiency boost!
  - Meet partner now

# Size and Height of Binary Trees

- ▶ Notation:
  - Let  $T$  be a tree
  - Write  $h(T)$  for the height of the tree, and
  - $N(T)$  for the size (i.e., number of nodes) of the tree
- ▶ Given  $h(T)$ , what are the bounds on  $N(T)$ ?
  - $N(T) \leq \text{-----}$  and  $N(T) \geq \text{-----}$
- ▶ Given  $N(T)$ , what are the bounds on  $h(T)$ ?
  - Solve each inequality for  $h(T)$  and combine

# Extreme Trees

- ▶ A tree with the maximum number of nodes for its height is a **full tree**.
  - Its height is  **$O(\log N)$**
- ▶ A tree with the minimum number of nodes for its height is essentially a \_\_\_\_\_
  - Its height is  **$O(N)$**
- ▶ Height matters!
  - Recall that the algorithms for search, insertion, and deletion in a binary search tree are  **$O(h(T))$**

To prove recursive properties (on trees), we use a technique called mathematical induction

- ▶ Actually, we use a variant called *strong induction* :



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California

# Strong Induction

- ▶ To prove that  $p(n)$  is true for all  $n \geq n_0$ :
  - Prove that  $p(n_0)$  is true (base case), and
  - For all  $k > n_0$ , prove that if we assume  $p(j)$  is true for  $n_0 \leq j < k$ , then  $p(k)$  is also true
- ▶ An analogy for those who took MA275:
  - Regular induction uses the previous domino to knock down the next
  - Strong induction uses all the previous dominos to knock down the next!
- ▶ Warmup: prove the arithmetic series formula
- ▶ Actual: prove the formula for  $N(T)$

# Exam Review

# The Big Picture

- ▶ All data structures really boil down to:
  - Continuous memory (arrays), or
  - Nodes and pointers (**linked lists, trees, graphs**)
- ▶ Let's draw pics of each
- ▶ Then you do the questions on the back with a partner as exam review
  
- ▶ Then time for questions